

The Open University of Sri Lanka
B.Sc. Degree Programme – Level 05
Open Book Test (OBT) - 2009/2010
Pure Mathematics / Computer Science
PMU 3294/PME 5294/CSU 3276 – Discrete Mathematics



Model Answer

01. (a) The given statement is : “If two triangles are similar then they are congruent”.

Converse: “If two triangles are congruent then they are similar”,

Inverse: “If two triangles are not similar then they are not congruent”,

Contra - positive: “If two triangles are not congruent then they are not similar”.

(b) Let p and q be two statements. Then consider the following truth table:

p	q	$p \Leftrightarrow q$	$\neg p$	$\neg q$	$(\neg p \vee q)$	$(p \vee \neg q)$	$(\neg p \vee q) \wedge (p \vee \neg q)$
T	T	T	F	F	T	T	T
T	F	F	F	T	F	T	F
F	T	F	T	F	T	F	F
F	F	T	T	T	T	T	T

Truth values in the column under the heading $(p \Leftrightarrow q)$ are the same as truth values in the column under the heading $(\neg p \vee q) \wedge (p \vee \neg q)$.

Hence $(p \Leftrightarrow q)$ is logically equivalent to $(\neg p \vee q) \wedge (p \vee \neg q)$.

(c) (i) “ $\forall x \in \mathbb{R}, x^2 > 0$ ” is a false statement.

Since $x = 0 (\in \mathbb{R})$, but $x^2 = 0$ is not greater than 0.

(ii) “ $(1^3 + 2^3 + 3^3 + \dots + n^3) = (1 + 2 + 3 + \dots + n)^2$ ” is a true statement.

We will prove this by method of induction

We know that $(1^3 + 2^3 + 3^3 + \dots + n^3) = (1 + 2 + 3 + \dots + n)^2 = \left[\frac{n(n+1)}{2} \right]^2$.

When $n = 1$:

$1^3 = 1^2$ is true, therefore the result is true for $n = 1$

Assume that the result is true for $n = k$:

$$\text{i.e. } (1^3 + 2^3 + 3^3 + \dots + k^3) = (1 + 2 + 3 + \dots + k)^2 = \left[\frac{k(k+1)}{2} \right]^2$$

Now consider the result when $n = k + 1$:

$$\begin{aligned} (1^3 + 2^3 + 3^3 + \dots + k^3 + (k+1)^3) &= \left[\frac{k(k+1)}{2} \right]^2 + (k+1)^3 \\ &= \frac{k^4 + 2k^3 + k^2}{4} + k^3 + 3k^2 + 3k + 1 \\ &= \frac{k^4 + 6k^3 + 13k^2 + 12k + 4}{4} \\ &= \left[\frac{(k+1)(k+2)}{2} \right]^2 \\ &= [1 + 2 + 3 + \dots + k + (k+1)]^2 \end{aligned}$$

So the result is true for $n = k + 1$ also

Therefore by principle of mathematical induction the result is true $\forall n \in \mathbb{N}$.

(iii) " $\forall x \in \mathbb{Q}, x^2 \in \mathbb{Q}$ " is a true statement.

If $x \in \mathbb{Q}$, then $x = \frac{p}{q}$ where $p, q \in \mathbb{Z}$ with $q \neq 0$

$$\text{So } x^2 = \frac{p^2}{q^2} \in \mathbb{Q} \quad (\because p, q \in \mathbb{Z} \text{ therefore } p^2, q^2 \in \mathbb{Z} \text{ } q^2 \neq 0)$$

In Addition, suppose if the question: 'is the statement

" $\forall x \in \mathbb{R} \setminus \mathbb{Q}, x^2 \in \mathbb{R} \setminus \mathbb{Q}$ " true or false?'

Then it is a false statement.

$$\text{Since } \sqrt{2} \in \mathbb{R} \setminus \mathbb{Q} \quad \text{but } (\sqrt{2})^2 = 2 \notin \mathbb{R} \setminus \mathbb{Q}$$

02. (a) (i) $R = \{(x, y) \mid (x \in \mathbb{R}) \wedge (y \in \mathbb{R}) \wedge (y = x^2)\}$ is not a partial order

Since R is not reflexive (\because there is no $_2 R_2$ i.e. $2 \neq 2^2$)

(ii) $R = \{(x, y) \mid (x \in A) \wedge (y \in A) \wedge ("x \text{ divides } y")\},$

where $A = \{1, 2, 3, 4, 6\}$. Here

$$R = \left\{ \begin{array}{l} (1, 1), (1, 2), (1, 3), (1, 4), (1, 6), (2, 2), \\ (2, 4), (2, 6), (3, 3), (3, 6), (4, 4), (6, 6) \end{array} \right\}$$

and R is reflexive, anti-symmetric and transitive,
therefore R is a partial order.

(iii) $R = \{(x, y) \mid (x \in \mathbb{R}) \wedge (y \in \mathbb{R}) \wedge (y \leq x)\}$ is a partial order

$_x R_x \forall x \in \mathbb{R}$ i.e. $x \leq x \forall x \in \mathbb{R}$, so R is reflexive

If $_x R_y$ and $_y R_x$, then $x \leq y$ and $y \leq x$ this gives $x = y$

So R is anti-symmetric

If $_x R_y$ and $_y R_z$, then $x \leq y$ and $y \leq z$ this gives $x \leq z$

So R is transitive and therefore R is a partial order.

(b) Let f be a homomorphism from a group $(G, *)$ into a group $(G', *')$.

Let e, e' be the identity elements of G and G' respectively.

Then $f(e) = e'$,

Let $x \in G$ and suppose that x^{-1} is the inverse of x in G .

Then $f(x * x^{-1}) = f(x) *' f(x^{-1})$. Also

$$x * x^{-1} = e$$

$$f(x * x^{-1}) = f(e)$$

$$f(x) *' f(x^{-1}) = e' \quad (\because f \text{ is a homomorphism})$$

This gives $f(x^{-1})$ is the inverse of $f(x)$

$$\text{i.e. } [f(x)]^{-1} = f(x^{-1}).$$



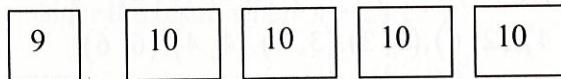
(c) The operation “*subtraction*” in \mathbb{Z} , the set of integers, is not associative:

Since

$$\mathbb{Z} = \{\dots -2, -1, 0, 1, 2, \dots\}$$

$$(12 - 6) - 2 = 6 - 2 = 4 \quad \text{but} \quad 12 - (6 - 2) = 12 - 4 = 8$$

03. (a) Digits are 0,1,2,3,4,5,6,7,8,9 and number of digits is 10



So the number of telephone numbers can be assigned is 9×10^4 .

(b) Let the number of r - permutations of n objects is denoted by ${}^n P_r$
and the number of r - objects subsets of n - objects set is denoted by ${}^n C_r$.

$$\text{Number of } r\text{-subsets of an } n\text{-objects set} = {}^n C_r = \frac{n!}{r!(n-r)!}$$

$$\text{Number of } r\text{-permutations of } r\text{ objects} = r!$$

Then by the product principle, the number of r - permutations of n objects is

$$r! {}^n C_r = r! \frac{n!}{r!(n-r)!} = \frac{n!}{(n-r)!} = {}^n P_r.$$

(c) There are 20 Mathematics students and 15 Computer Science students in the Discrete Mathematics class and the number of ways of selecting **12 students** from the class, if

- (i) there are no restrictions $= {}^{35} C_{12}$,
- (ii) all must be mathematics major $= {}^{20} C_{12} {}^{15} C_0$,
- (iii) all must belong to the same discipline $= {}^{20} C_0 {}^{15} C_{12} + {}^{20} C_{12} {}^{15} C_0$,
- (iv) the two disciplines must have the same number of representatives $= {}^{20} C_6 {}^{15} C_6$.

